



# DIAGRAMMATIC QUANTUM CIRCUIT COMPRESSION FOR HAMILTONIAN SIMULATION

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Computing Analytics (PGI-12)

# CONTRIBUTION

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# OUTLINE

- Motivation
- ZX-calculus and Pauli-gadgets
- Compression of Quantum Circuits for Hamiltonian Simulation

# MOTIVATION

- NISQ: Restrictions on number of qubits and gates
- Shallow and simple quantum circuit structures are crucial
- Increasing popularity of diagrammatic approaches to quantum computing
- Used tool: ZX-calculus

# ZX-calculus and Pauli-gadgets

# ZX-DIAGRAMS: GENERATORS

Spiders:

$$m \left\{ \begin{array}{c} \text{---} \\ \vdots \\ \text{---} \end{array} \right\} \alpha \left\{ \begin{array}{c} \text{---} \\ \vdots \\ \text{---} \end{array} \right\} n := \underbrace{|0, \dots, 0\rangle}_n \underbrace{\langle 0, \dots, 0|}_m + e^{i\alpha} \underbrace{|1, \dots, 1\rangle}_n \underbrace{\langle 1, \dots, 1|}_m, \quad (1)$$

$$m \left\{ \begin{array}{c} \text{---} \\ \vdots \\ \text{---} \end{array} \right\} \alpha \left\{ \begin{array}{c} \text{---} \\ \vdots \\ \text{---} \end{array} \right\} n := \underbrace{|+, \dots, +\rangle}_n \underbrace{\langle +, \dots, +|}_m + e^{i\alpha} \underbrace{|-, \dots, -\rangle}_n \underbrace{\langle -, \dots, -|}_m. \quad (2)$$

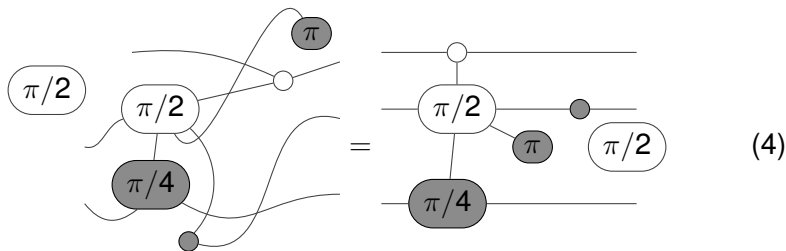
Hadamard:

$$\text{---} \square \text{---} = 1/\sqrt{2} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad (3)$$

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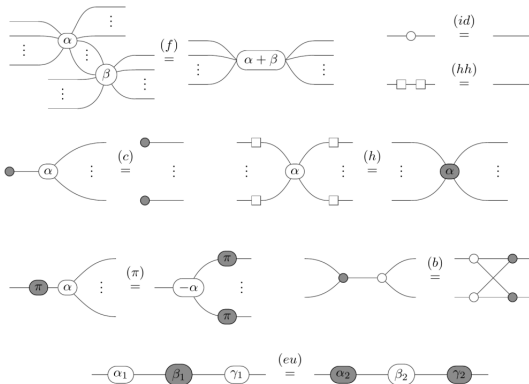
<sup>1</sup>Van de Wetering, J.: ZX-calculus for the working quantum computer scientist, 2020, arXiv: 2012.13966 [quant-ph].

# ZX-DIAGRAMS: ONLY CONNECTIVITY MATTERS



<sup>1</sup>Van de Wetering, J.: ZX-calculus for the working quantum computer scientist, 2020, arXiv: 2012.13966 [quant-ph].

# ZX-CALCULUS



A presentation of the ZX-calculus. All rules hold for all  $\alpha, \beta \in \mathbb{R}$  and due to (h) and (hh) for all colors interchanged. These rules only hold up to a non-zero scalar.

<sup>1</sup>Van de Wetering, J.: ZX-calculus for the working quantum computer scientist, 2020, arXiv: 2012.13966 [quant-ph].



# PHASE-GADGETS

The Z-phase gadgets  $\Phi_n^Z(\alpha) : \mathbb{C}^{\otimes n} \longrightarrow \mathbb{C}^{\otimes n}$  are a family of unitary maps and are recursively defined as:

$$\begin{aligned}\Phi_1^Z(\alpha) &:= Z(\alpha), \\ \Phi_{n+1}^Z(\alpha) &:= (\mathbb{1}_n \otimes \text{CNOT}(n+1, n))(\Phi_n^Z \otimes \mathbb{1}_1)(\mathbb{1}_n \otimes \text{CNOT}(n+1, n))\end{aligned}\quad (5)$$

In ZX-calculus:

$$\Phi_n^Z(\alpha) = \text{Diagram 1}, \quad \Phi_n^X(\alpha) = \text{Diagram 2}\quad (6)$$

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<sup>2</sup>Cowtan, A. et al.: Phase Gadget Synthesis for Shallow Circuits. Electronic Proceedings in Theoretical Computer Science 318, pp. 213–228, 2020, <https://doi.org/10.4204%2Feptcs.318.13>.

# PAULI-GADGETS

Let  $s$  be a word over the alphabet  $\{X, Y, Z\}$ , then a Pauli-gadget  $P(\alpha, s)$  is defined as  $U(s)\Phi_{|s|}^Z(\alpha)U(s)^\dagger$  where the unitary  $U(s)$  is defined recursively over  $s$ :

$$U(Zs') = I \otimes U(s'), \quad U(Ys') = X(\pi/2) \otimes U(s'), \quad U(Xs') = H \otimes U(s') \quad (7)$$

Example:

The diagram shows the decomposition of the Pauli-gadget  $e^{-i\alpha ZXIY}$  into two equivalent quantum circuits. The left circuit consists of a sequence of gates: a CNOT from the top qubit to the second, a CNOT from the second to the third, and a CNOT from the third to the top. This is followed by a rotation gate  $\frac{\pi}{2}$  on the top qubit, a CNOT from the top to the second, a rotation gate  $-\frac{\pi}{2}$  on the top qubit, and a CNOT from the top to the second. The right circuit consists of a sequence of gates: a CNOT from the top to the second, a CNOT from the second to the third, and a CNOT from the third to the top, followed by a rotation gate  $2\alpha$  on the top qubit. The two circuits are shown to be equivalent with a double-line arrow between them.

$$e^{-i\alpha ZXIY} = \text{[Circuit 1]} := \text{[Circuit 2]} \quad (8)$$

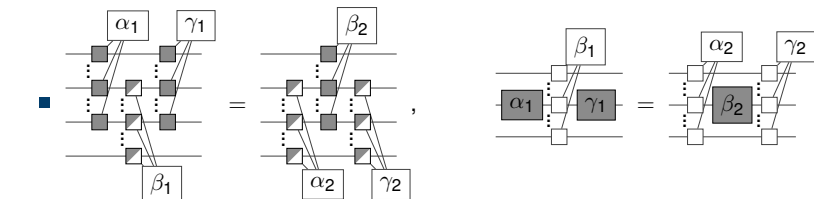
<sup>2</sup>Cowan, A. et al.: Phase Gadget Synthesis for Shallow Circuits. Electronic Proceedings in Theoretical Computer Science 318, pp. 213–228, 2020, <https://doi.org/10.4204%2Feptcs.318.13>.

# FUSION AND TURNOVER RULES

- Let  $s$  be a Pauli-string, then for all  $\alpha$  and  $\beta$ :

$$P(\alpha, s)P(\beta, s) = P(\alpha + \beta, s) \quad (9)$$

- Two gadgets with an even number of legs on the same wires commute.



<sup>2</sup>Cowan, A. et al.: Phase Gadget Synthesis for Shallow Circuits. Electronic Proceedings in Theoretical Computer Science 318, pp. 213–228, 2020, <https://doi.org/10.4204/2Feptcs.318.13>.

<sup>3</sup>Winderl, D. et al.: A recursively partitioned approach to architecture-aware ZX Polynomial synthesis and optimization, 2023, arXiv: 2303.17366 [quant-ph]

# Diagrammatic Compression for Hamiltonian Simulation

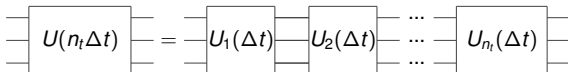
# HAMILTONIAN SIMULATION

Time evolution operator:

$$U(t_0, t_1) = \tau \exp \left( -i \int_{t_0}^{t_1} H(t) dt \right), \quad H(t) = \sum_{\ell} H_{\ell}(t) \quad (10)$$

By discretization in time and Trotter decomposition:  
Approximate time evolution operator by

$$U(n_t \Delta t) = \prod_{k=0}^{n_t-1} U_{n_t-k}(\Delta t), \quad U_k(\Delta t) = \prod_{\ell} \exp(-iH_{\ell,k} \Delta t). \quad (11)$$



<sup>5</sup>Camps, D. et al.: An Algebraic Quantum Circuit Compression Algorithm for Hamiltonian Simulation. SIAM Journal on Matrix Analysis and Applications 43 (3), pp. 1084–1108, 2022, <https://doi.org/10.1137%2F21m1439298>.

# COMPRESSION USING TURNOVER OPERATIONS

## Definition

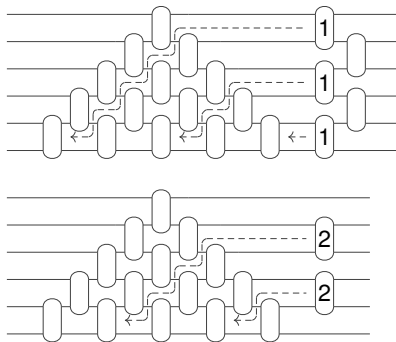
Define a *block*  $B_i(\vec{\Theta})$  as a structure that satisfies

- 1 Fusion:  $B_i(\vec{\alpha})B_i(\vec{\beta}) = B_i(\vec{a})$
- 2 Commutation:  $B_i(\vec{\alpha})B_j(\vec{\beta}) = B_j(\vec{\beta})B_i(\vec{\alpha})$  for  $|i - j| > 1$
- 3 Turnover:  $B_i(\vec{\alpha})B_{i+1}(\vec{\beta})B_i(\vec{\gamma}) = B_{i+1}(\vec{a})B_i(\vec{b})B_{i+1}(\vec{c})$

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<sup>4</sup>Kökcü, E. et al.: Algebraic compression of quantum circuits for Hamiltonian evolution. Physical Review A 105 (3), 2022, <https://doi.org/10.1103/PhysRevA.105.032420>.

# MERGING A TROTTER STEP INTO A CONSTANT-DEPTH TRIANGLE



<sup>5</sup>Camps, D. et al.: An Algebraic Quantum Circuit Compression Algorithm for Hamiltonian Simulation. SIAM Journal on Matrix Analysis and Applications 43 (3), pp. 1084–1108, 2022, <https://doi.org/10.1137%F21m1439298>.

# COMPRESSION ALGORITHM

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**Input:** Trotter circuit  $C$  on  $N$  spins with  $n_t$  time-steps,  $n_t > N/2$

**Output:** Compressed  $N \times N$  Trotter circuit  $C'$  equivalent to  $C$

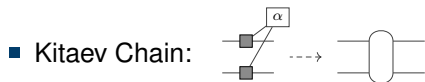
- 1:  $C' \leftarrow \text{TriangleCircuit}(C'[:, 1 : N])$
  - 2: **for**  $l = N + 1$  to  $2n_t$  **do**
  - 3:      $\text{MergeLayer}(C', C[:, l])$
  - 4:  $C' \leftarrow \text{SquareCircuit}(C')$
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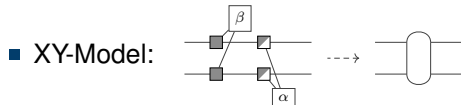
<sup>4</sup>Camps, D. et al.: An Algebraic Quantum Circuit Compression Algorithm for Hamiltonian Simulation. SIAM Journal on Matrix Analysis and Applications 43 (3), pp. 1084–1108, 2022, <https://doi.org/10.1137/F21m1439298>.



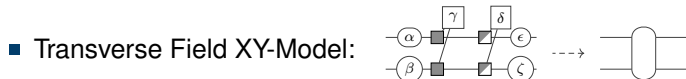
# BLOCKS IN ZX-CALCULUS



$$H(t) = \sum_{i=1}^{N-1} J_i^{\alpha_i} (t) \sigma_i^{\alpha_i} \sigma_{i+1}^{\alpha_i}, \quad \text{for } \alpha_i \in \{x, y, z\}, \alpha_i \neq \alpha_{i+1} \quad (12)$$

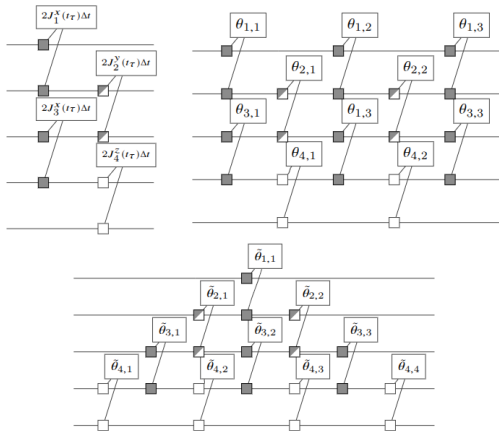


$$H(t) = \sum_{i=1}^{N-1} J_i^{\alpha} (t) \sigma_i^{\alpha} \sigma_{i+1}^{\alpha} + J_i^{\beta} (t) \sigma_i^{\beta} \sigma_{i+1}^{\beta}, \quad \text{for } \alpha \neq \beta \in \{x, y, z\} \quad (13)$$



$$H(t) = \sum_{i=1}^{N-1} J_i^x (t) \sigma_i^x \sigma_{i+1}^x + J_i^y (t) \sigma_i^y \sigma_{i+1}^y + \sum_{i=1}^N h_i^z (t) \sigma_i^z \quad (14)$$

# KITAEV-COMPRESSION FOR 5 QUBITS



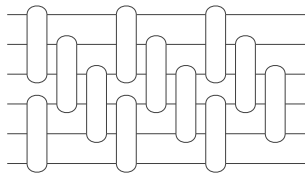
One Trotter timestep (top left). The constant-depth square circuit for 5 qubits (top right).

The constant-depth triangle circuit for 5 qubits (bottom).

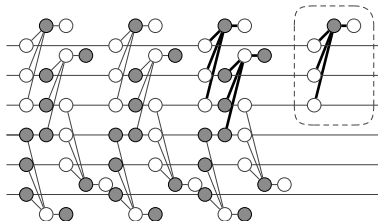
# KITAEV CHAIN WITH THREE-BODY INTERACTION

$$H(t) = \sum_{i=1}^{N-2} J_i^{\alpha_i}(t) \sigma_i^{\alpha_i} \sigma_{i+1}^{\alpha_i} \sigma_{i+2}^{\alpha_i}, \quad (15)$$

for  $\alpha_i \neq \alpha_{i+1}$  and the restriction that  $\alpha_i \in \{x, z\}$ ,  $\alpha_i \in \{x, y\}$  or  $\alpha_i \in \{y, z\}$  for all  $i$  (i.e. only two different types of Pauli- $\alpha$  occur).



The constant depth square for 6 qubits.



The compression of one Pauli-ZZZ gadget (highlighted) in ZX-calculus for 6 qubits.

# SUMMARY AND OUTLOOK

- Usefulness of diagramming tools (ZX-calculus) to show circuit compression
- Re-derived constant-depth properties for several (Ising-, Kitaev-, XY-, TFXY-, TFIM-) models
- Showed constant-depth properties for three-spin interaction Kitaev model
- Indications for constant-depth properties for three-spin transverse-field Ising model

Thanks for your attention!

# Hamiltonian Simulation

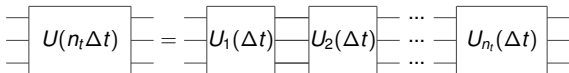
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Time evolution operator:

$$U(t_0, t_1) = \tau \exp \left( -i \int_{t_0}^{t_1} H(t) dt \right), \quad H(t) = \sum_{\ell} H_{\ell}(t) \quad (16)$$

By discretization in time and Trotter decomposition:  
Approximate time evolution operator by

$$U(n_t \Delta t) = \prod_{k=0}^{n_t-1} U_{n_t-k}(\Delta t), \quad U_k(\Delta t) = \prod_{\ell} \exp(-iH_{\ell,k} \Delta t). \quad (17)$$

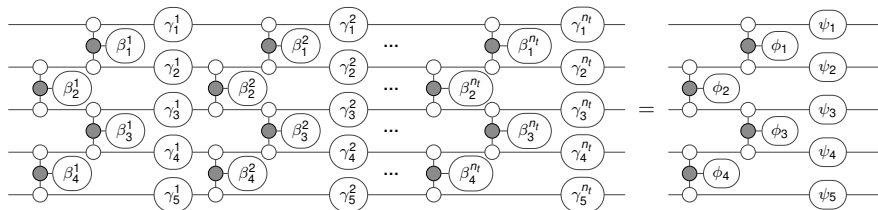


<sup>5</sup>Camps, D. et al.: An Algebraic Quantum Circuit Compression Algorithm for Hamiltonian Simulation. SIAM Journal on Matrix Analysis and Applications 43 (3), pp. 1084–1108, 2022, <https://doi.org/10.1137%2F21m1439298>.

# CLASSICAL ISING MODEL

$$H(t) = \sum_{i=1}^{n-1} J_i^\alpha(t) \sigma_i^\alpha \sigma_{i+1}^\alpha + \sum_{i=1}^n h_i^\alpha(t) \sigma_i^\alpha, \quad \alpha \in \{x, y, z\} \quad (18)$$

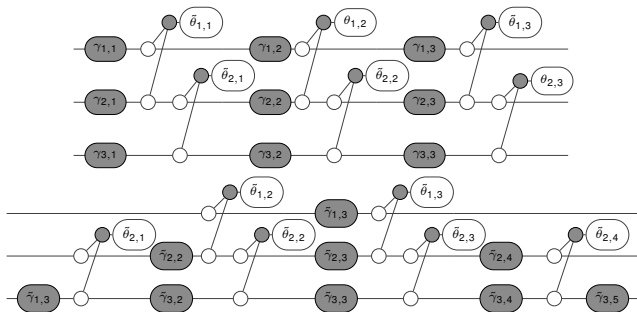
Approximate time evolution operator for  $\alpha = z$ :





# TRANSVERSE FIELD ISING MODEL

$$H(t) = \sum_{i=1}^{N-1} J_i^\alpha(t) \sigma_i^\alpha \sigma_{i+1}^\alpha + \sum_{i=1}^N h_i^\beta(t) \sigma_i^\beta \quad (19)$$



# TRANSVERSE FIELD ISING MODEL WITH THREE-BODY INTERACTION

$$H(t) = \sum_{i=1}^{N-2} J_i^z(t) \sigma_i^z \sigma_{i+1}^z \sigma_{i+2}^z + \sum_{i=1}^N h_i^x(t) \sigma_i^x \quad (20)$$

